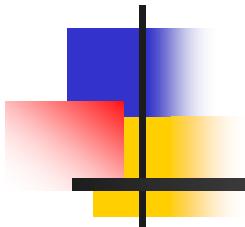


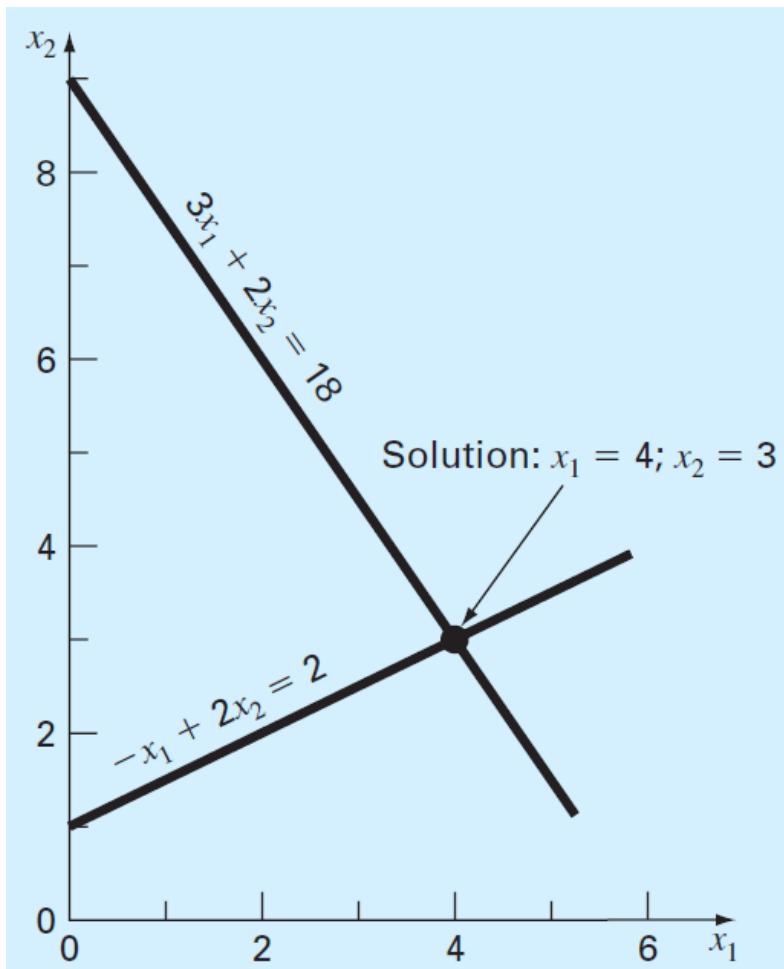
# Gauss Elimination



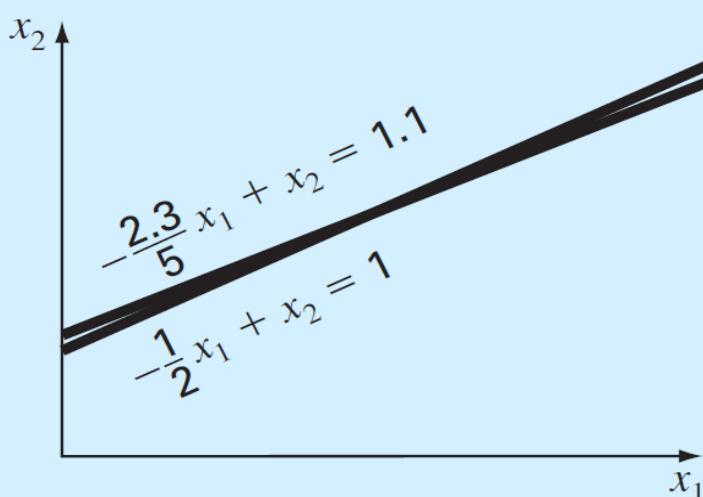
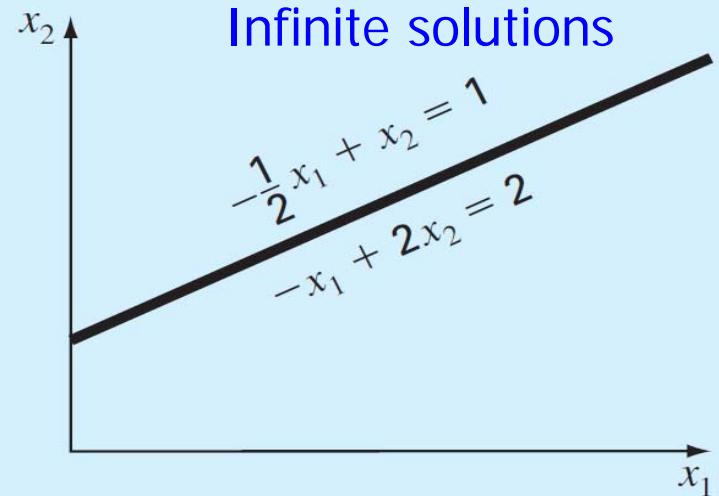
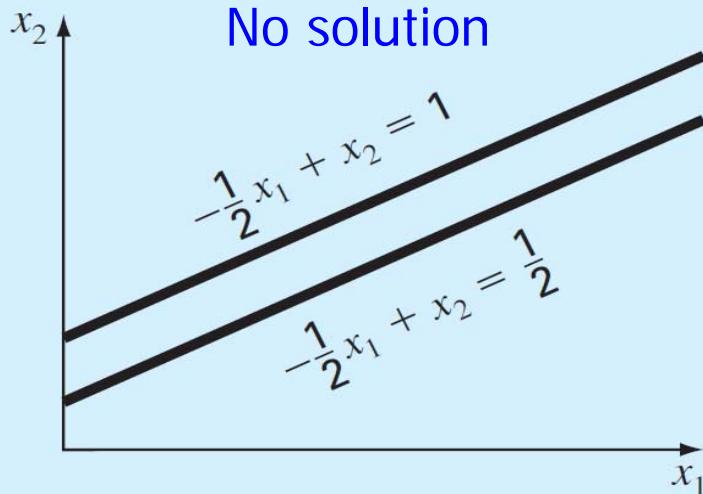
Hsiao-Lung Chan  
Dept Electrical Engineering  
Chang Gung University, Taiwan  
[chanhl@mail.cgu.edu.tw](mailto:chanhl@mail.cgu.edu.tw)

# Solving small numbers of equations by graphical method

- The location of the intercept provides a solution



# Singular and ill-conditioned systems



III-conditioned system where  
the slopes are so close that the  
point of intersection is difficult  
to detect visually

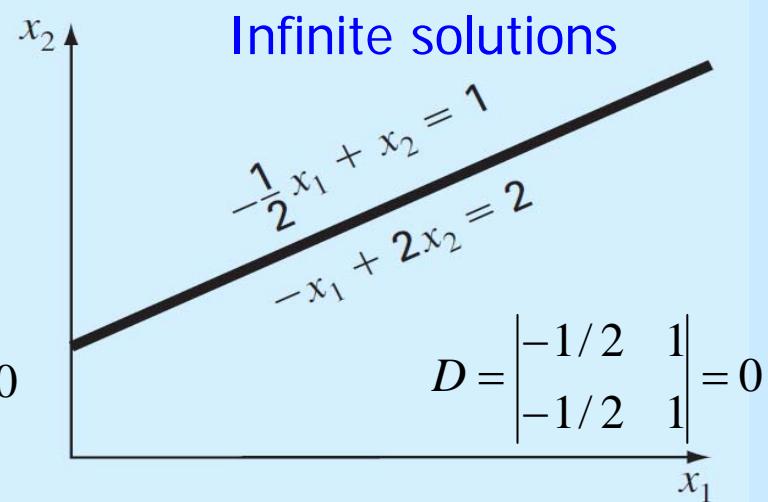
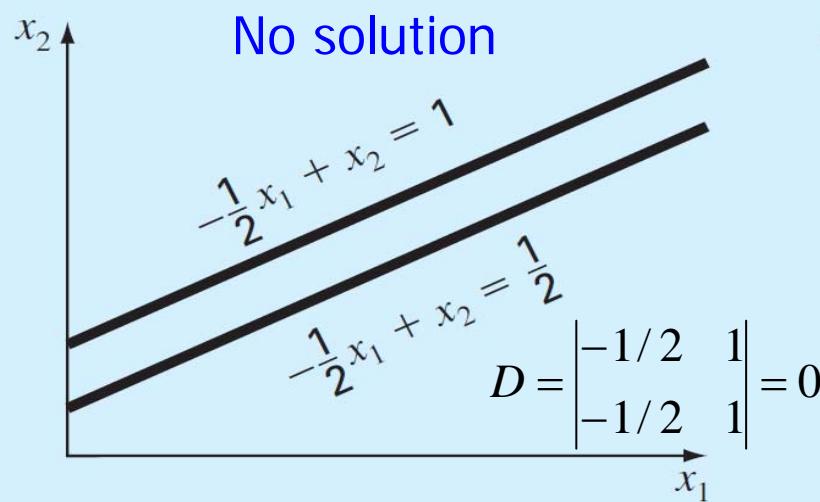
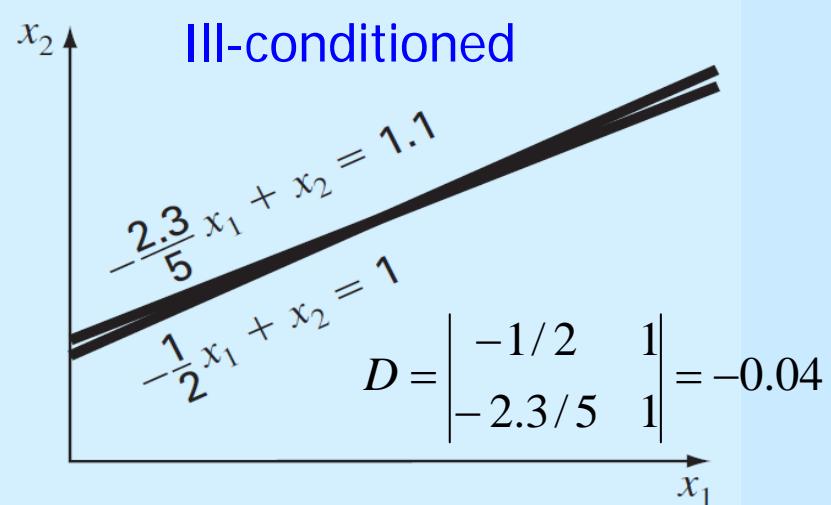
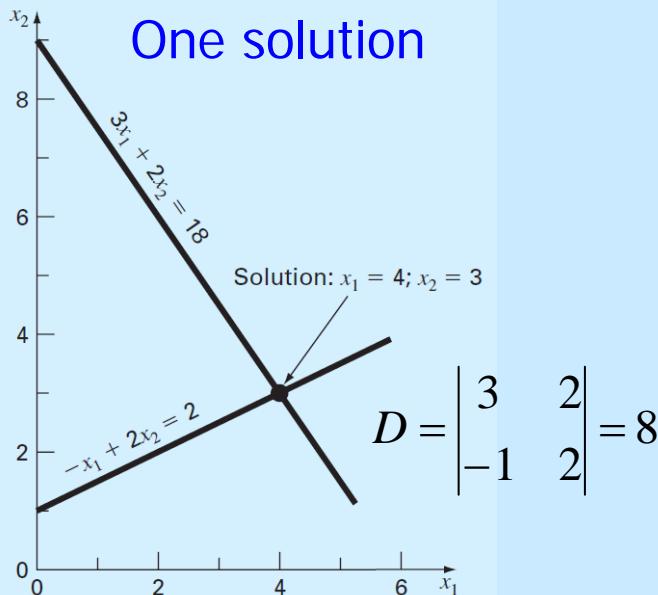
# Determinants

---

$$D = |A|$$

$$\begin{array}{c} 1 \times 1 \quad |a_{11}| = a_{11} \\ \hline 2 \times 2 \quad \begin{vmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{vmatrix} = a_{11}a_{22} - a_{12}a_{21} \\ \hline 3 \times 3 \quad \begin{vmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{vmatrix} = a_{11} \begin{vmatrix} a_{22} & a_{23} \\ a_{32} & a_{33} \end{vmatrix} - a_{12} \begin{vmatrix} a_{21} & a_{23} \\ a_{31} & a_{33} \end{vmatrix} + a_{13} \begin{vmatrix} a_{21} & a_{22} \\ a_{31} & a_{32} \end{vmatrix} \end{array}$$

# Determinants in linear equations



## Cramer's Rule

---

- Each unknown in a system of linear algebraic equations may be expressed as a fraction of two determinants

$$x_1 = \frac{\begin{vmatrix} b_1 & a_{12} & a_{13} \\ b_2 & a_{22} & a_{23} \\ b_3 & a_{32} & a_{33} \end{vmatrix}}{D} \quad \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix}$$

A=[0.3 0.52 1; 0.5 1 1.9; 0.1 0.3 0.5];

D=det(A);

A(:,1)=[-0.01; 0.67; -0.44];

x1=det(A)/D;

# Naive Gauss elimination

- A sequential process of removing unknowns from equations
- 'Naive' means the process does not check for division-by-zero

$$\left[ \begin{array}{ccc|c} a_{11} & a_{12} & a_{13} & b_1 \\ a_{21} & a_{22} & a_{23} & b_2 \\ a_{31} & a_{32} & a_{33} & b_3 \end{array} \right] \quad \left. \begin{array}{l} \\ \\ \end{array} \right\} (a) \text{ Forward elimination}$$

↓

$$\left[ \begin{array}{ccc|c} a_{11} & a_{12} & a_{13} & b_1 \\ a'_{22} & a'_{23} & & b'_2 \\ a''_{33} & & & b''_3 \end{array} \right] \quad \left. \begin{array}{l} \\ \\ \end{array} \right\} (b) \text{ Back substitution}$$

↓

$$x_3 = b''_3 / a''_{33}$$
$$x_2 = (b'_2 - a'_{23}x_3) / a'_{22}$$
$$x_1 = (b_1 - a_{13}x_3 - a_{12}x_2) / a_{11}$$

# Forward elimination of unknown

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + \cdots + a_{1n}x_n = b_1$$

$$a_{21}x_1 + a_{22}x_2 + a_{23}x_3 + \cdots + a_{2n}x_n = b_2$$

$\vdots$        $\vdots$

$$a_{n1}x_1 + a_{n2}x_2 + a_{n3}x_3 + \cdots + a_{nn}x_n = b_n$$

multiplied by  $a_{21}/a_{11}$

$$a_{21}x_1 + \frac{a_{21}}{a_{11}}a_{12}x_2 + \frac{a_{21}}{a_{11}}a_{13}x_3 + \cdots + \frac{a_{21}}{a_{11}}a_{1n}x_n = \frac{a_{21}}{a_{11}}b_1$$

$$\xrightarrow{\quad\quad\quad} \left( a_{22} - \frac{a_{21}}{a_{11}}a_{12} \right)x_2 + \cdots + \left( a_{2n} - \frac{a_{21}}{a_{11}}a_{1n} \right)x_n = b_2 - \frac{a_{21}}{a_{11}}b_1$$

subtracted from  
2<sup>nd</sup> equation

$$a'_{22}x_2 + \cdots + a'_{2n}x_n = b'_2$$

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + \cdots + a_{1n}x_n = b_1$$

continuing n-2  
eliminations



$$a'_{22}x_2 + a'_{23}x_3 + \cdots + a'_{2n}x_n = b'_2$$

$$a''_{33}x_3 + \cdots + a''_{3n}x_n = b''_3$$

$\ddots$

$$a^{(n-1)}_{nn}x_n = b_n^{(n-1)}$$

## backward substitution

---

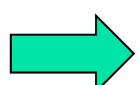
$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + \cdots + a_{1n}x_n = b_1$$

$$a'_{22}x_2 + a'_{23}x_3 + \cdots + a'_{2n}x_n = b'_2$$

$$a''_{33}x_3 + \cdots + a''_{3n}x_n = b''_3$$

⋮ ⋮

$$a_{nn}^{(n-1)}x_n = b_n^{(n-1)}$$



$$x_n = \frac{b_n^{(n-1)}}{a_{nn}^{(n-1)}}$$

$$x_i = \frac{b_i^{(i-1)} - \sum_{j=i+1}^n a_{ij}^{(i-1)}x_j}{a_{ii}^{(i-1)}} \quad \text{for } i = n-1, n-2, \dots, 1$$

# Example of Gauss elimination

## Forward elimination

$$\begin{array}{l} 3x_1 - 0.1x_2 - 0.2x_3 = 7.85 \\ 0.1x_1 + 7x_2 - 0.3x_3 = -19.3 \\ 0.3x_1 - 0.2x_2 + 10x_3 = 71.4 \end{array} \quad \begin{array}{l} f_{21}=0.1/3 \\ f_{31}=0.3/3 \\ f_{32}=-0.19/7.00333 \end{array}$$

↓

$$\begin{array}{l} 3x_1 - 0.1x_2 - 0.2x_3 = 7.85 \\ 7.00333x_2 - 0.293333x_3 = -19.5617 \\ - 0.190000x_2 + 10.0200x_3 = 70.6150 \end{array}$$
$$\begin{array}{l} 3x_1 - 0.1x_2 - 0.2x_3 = 7.85 \\ 7.00333x_2 - 0.293333x_3 = -19.5617 \\ 10.0120x_3 = 70.0843 \end{array}$$

## Backward substitution

$$x_3 = \frac{70.0843}{10.0120} = 7.00003 \quad x_2 = \frac{-19.5617 + 0.293333(7.00003)}{7.00333} = -2.50000$$

$$x_1 = \frac{7.85 + 0.1(-2.50000) + 0.2(7.00003)}{3} = 3.00000$$

## M-file to implement naive Gauss elimination

---

```
function x = GaussNaive(A,b)
[m,n] = size(A);
if m~=n, error('Matrix A must be square'); end
nb = n+1;
Aug = [A b];
% forward elimination
for k=1:n-1
    for i=k+1:n
        factor=Aug(i,k)/Aug(k,k);
        Aug(i,k:nb)=Aug(i,k:nb)-factor*Aug(k,k:nb);
    end
end
```

## GaussNaive M-file (cont.)

---

```
% back substitution  
x = zeros(n,1);  
x(n) = Aug(n,nb)/Aug(n,n);  
for i = n-1:-1:1  
    x(i) = (Aug(i,nb)-Aug(i,i+1:n)*x(i+1:n))/Aug(i,i);  
end
```

# Program efficiency

---

- Execution time depends on the amount of floating-point operations (flops).

$$\sum_{i=1}^m 1 = 1 + 1 + 1 + \cdots + 1 = m \quad \sum_{i=k}^m 1 = m - k + 1$$

$$\sum_{i=1}^m i = 1 + 2 + 3 + \cdots + m = \frac{m(m+1)}{2} = \frac{m^2}{2} + O(m)$$

$$\sum_{i=1}^m i^2 = 1^2 + 2^2 + 3^2 + \cdots + m^2 = \frac{m(m+1)(2m+1)}{6} = \frac{m^3}{3} + O(m^2)$$

# Program efficiency of Gauss elimination

---

## ■ Forward elimination

```
for k=1:n-1
    for i=k+1:n
        factor=Aug(i,k)/Aug(k,k);
        Aug(i,k:nb)=Aug(i,k:nb)-factor*Aug(k,k:nb);
    end
end
```

Outer Loop <i>k</i>	Inner Loop <i>i</i>	Addition/Subtraction Flops	Multiplication/Division Flops
1	2, <i>n</i>	$(n - 1)(n)$	$(n - 1)(n + 1)$
2	3, <i>n</i>	$(n - 2)(n - 1)$	$(n - 2)(n)$
:	:		
<i>k</i>	<i>k</i> + 1, <i>n</i>	$(n - k)(n + 1 - k)$	$(n - k)(n + 2 - k)$
:	:		
<i>n</i> - 1	<i>n</i> , <i>n</i>	(1)(2)	(1)(3)

---

## Program efficiency of Gauss elimination (cont.)

---

- Total addition/subtraction flops of forward elimination

$$\sum_{k=1}^{n-1} (n-k)(n+1-k) = \sum_{k=1}^{n-1} [n(n+1) - k(2n+1) + k^2]$$

$$n(n+1) \sum_{k=1}^{n-1} 1 - (2n+1) \sum_{k=1}^{n-1} k + \sum_{k=1}^{n-1} k^2$$

$$\rightarrow [n^3 + O(n^2)] - [n^3 + O(n^2)] + \left[ \frac{1}{3}n^3 + O(n^2) \right] = \frac{n^3}{3} + O(n^2)$$

- Similar analysis for multiplication/division flops

$$[n^3 + O(n^2)] - [n^3 + O(n)] + \left[ \frac{1}{3}n^3 + O(n^2) \right] = \frac{n^3}{3} + O(n^2)$$

- Total  $\frac{2n^3}{3} + O(n^2)$

## Program efficiency of Gauss elimination (cont.)

---

$$\underbrace{\frac{2n^3}{3} + O(n^2)}_{\text{Forward elimination}} + \underbrace{n^2 + O(n)}_{\text{Back substitution}} \xrightarrow{\text{as } n \text{ increases}} \frac{2n^3}{3} + O(n^2)$$

**TABLE 9.1** Number of flops for naive Gauss elimination.

<b><i>n</i></b>	<b>Elimination</b>	<b>Back Substitution</b>	<b>Total Flops</b>	<b><math>2n^3/3</math></b>	<b>Percent Due to Elimination</b>
10	705	100	805	667	87.58%
100	671550	10000	681550	666667	98.53%
1000	$6.67 \times 10^8$	$1 \times 10^6$	$6.68 \times 10^8$	$6.67 \times 10^8$	99.85%

## Problems arise with naive Gauss elimination

- If a coefficient along the diagonal is 0 (division by 0) or close to 0 (round-off error)

$$\begin{array}{l} 2x_2 + 3x_3 = 8 \\ 4x_1 + 6x_2 + 7x_3 = -3 \\ 2x_1 - 3x_2 + 6x_3 = 5 \end{array} \quad \begin{array}{l} 0.0003x_1 + 3.0000x_2 = 2.0001 \\ 1.0000x_1 + 1.0000x_2 = 1.0000 \\ -9999x_2 = -6666 \\ x_1 = \frac{2.0001 - 3(2/3)}{0.0003} \end{array}$$

$\downarrow$  Gauss elimination

Significant Figures	$x_2$	$x_1$	Absolute Value of Percent Relative Error for $x_1$
3	0.667	-3.33	1099
4	0.6667	0.0000	100
5	0.66667	0.30000	10
6	0.666667	0.330000	1
7	0.6666667	0.3330000	0.1

# Pivoting

---

- **Partial pivoting** : Determine the coefficient with the largest absolute value in the column below the pivot element. The rows can then be switched so that the largest element is the pivot element.
- **Complete pivoting** : If the rows to the right of the pivot element are also checked and columns switched.

## Example by partial pivoting

---

$$0.0003x_1 + 3.0000x_2 = 2.0001$$

$$1.0000x_1 + 1.0000x_2 = 1.0000$$

↓ Partial pivoting

$$1.0000x_1 + 1.0000x_2 = 1.0000$$

$$0.0003x_1 + 3.0000x_2 = 2.0001$$

↓ Gauss elimination

$$x_2 = 2/3$$

$$x_1 = \frac{1 - (2/3)}{1}$$

Significant Figures	$x_2$	$x_1$	Absolute Value of Percent Relative Error for $x_1$
3	0.667	0.333	0.1
4	0.6667	0.3333	0.01
5	0.66667	0.33333	0.001
6	0.666667	0.333333	0.0001
7	0.6666667	0.3333333	0.0000

## M-file to implement partial pivoting

---

```
function x = GaussPivot(A,b)
[m,n]=size(A);
if m~=n, error('Matrix A must be square'); end
nb=n+1;
Aug=[A b];

% forward elimination
for k = 1:n-1
    % partial pivoting
    [big,i]=max(abs(Aug(k:n,k)));
    ipr=i+k-1;
    if ipr~=k % rows exchange
        Aug([k,ipr],:)=Aug([ipr,k],:);
    end
end
```

## Partial pivoting M-file (cont.)

---

```
for i = k+1:n
    factor=Aug(i,k)/Aug(k,k);
    Aug(i,k:nb)=Aug(i,k:nb)-factor*Aug(k,k:nb);
end
end

% back substitution
x=zeros(n,1);
x(n)=Aug(n,nb)/Aug(n,n);
for i = n-1:-1:1
    x(i)=(Aug(i,nb)-Aug(i,i+1:n)*x(i+1:n))/Aug(i,i);
end
```

# Tridiagonal systems

---

- A banded system with a bandwidth of 3

$$\begin{bmatrix} f_1 & g_1 & & & \\ e_2 & f_2 & g_2 & & \\ e_3 & f_3 & g_3 & & \\ \cdot & \cdot & \cdot & \ddots & \\ & \cdot & \cdot & \cdot & \ddots & \cdot \\ & & \cdot & \cdot & \cdot & \cdot \\ e_{n-1} & f_{n-1} & g_{n-1} & & x_{n-1} \\ e_n & f_n & & & x_n \end{bmatrix} = \begin{bmatrix} r_1 \\ r_2 \\ r_3 \\ \cdot \\ \cdot \\ \cdot \\ r_{n-1} \\ r_n \end{bmatrix}$$

- Solved using the same method as Gauss elimination but with much less effort

## M-file for tridiagonal system solver

---

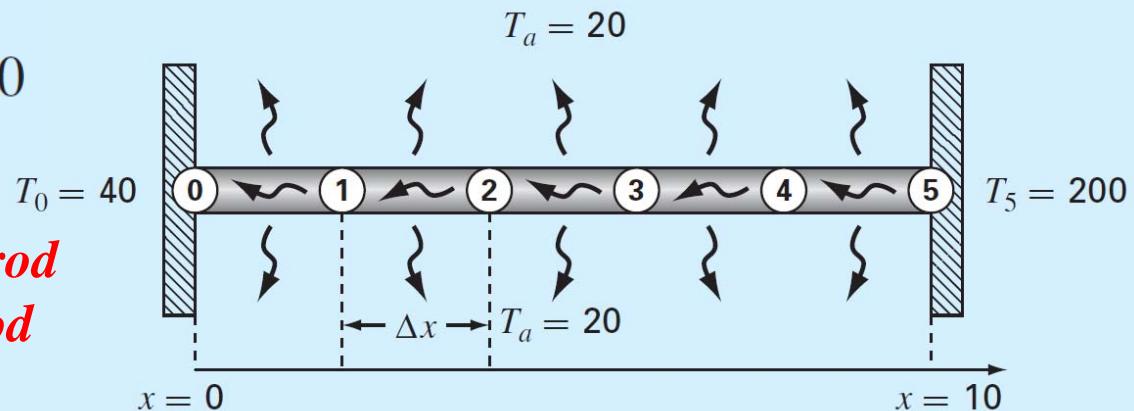
```
function x = Tridiag(e,f,g,r)
% e, f, g = subdiagonal, diagonal & superdiagonal vectors
n=length(f);
% forward elimination
for k=2:n
    factor = e(k)/f(k-1);
    f(k) = f(k) - factor*g(k-1);
    r(k) = r(k) - factor*r(k-1);
end
% back substitution
x(n) = r(n)/f(n);
for k = n-1:-1:1
    x(k) = (r(k)-g(k)*x(k+1))/f(k);
end
```

# A heated rod

- Steady-state, differential equation based on heat conservation

$$\frac{d^2T}{dx^2} + h'(T_a - T) = 0$$

*Heat flows through the rod  
as well as between the rod  
and the surrounding air*



- Finite-difference approximation

$$\frac{T_{i+1} - 2T_i + T_{i-1}}{\Delta x^2} + h'(T_a - T_i) = 0$$

## A heated rod (cont.)

Given  $h' = 0.01$ ,  $T_a = 20$ ,  $T(0) = 40$ ,  $T(10) = 200$ , get a solution

$$\frac{T_{i+1} - 2T_i + T_{i-1}}{\Delta x^2} + h'(T_a - T_i) = 0$$

Using  $\Delta x = 2$

$$-T_{i-1} + 2.04T_i - T_{i+1} = 0.8$$

$$\begin{aligned} -T_0 + 2.04T_1 - T_2 &= 0.8 \\ -T_1 + 2.04T_2 - T_3 &= 0.8 \\ -T_2 + 2.04T_3 - T_4 &= 0.8 \\ -T_3 + 2.04T_4 - T_5 &= 0.8 \end{aligned}$$

$$\begin{array}{l} T_0 = 40 \\ T_5 = 200 \end{array} \rightarrow$$

$$\begin{bmatrix} 2.04 & -1 & 0 & 0 \\ -1 & 2.04 & -1 & 0 \\ 0 & -1 & 2.04 & -1 \\ 0 & 0 & -1 & 2.04 \end{bmatrix} \begin{Bmatrix} T_1 \\ T_2 \\ T_3 \\ T_4 \end{Bmatrix} = \begin{Bmatrix} 40.8 \\ 0.8 \\ 0.8 \\ 200.8 \end{Bmatrix}$$

## A heated rod (cont.)

---

% Using left division

```
A=[2.04 -1 0 0; -1 2.04 -1 0; 0 -1 2.04 -1; 0 0 -1 2.04];  
b=[40.8 0.8 0.8 200.8]';  
T=(A\b)';
```

% Using tridiagonal system solver

```
e=[0 -1 -1 -1];  
f=[2.04 2.04 2.04 2.04];  
g=[-1 -1 -1 0];  
r=[40.8 0.8 0.8 200.8];  
T=Tridiag(e,f,g,r);
```

## Reference

---

- Steven C. Chapra "Applied Numerical Methods with MATLAB", 3rd ed., McGraw Hill, 2012.